

# Announcements

- Last homework: **HW11** on computability due May 1st

**Section plans:** this week : quiz + computability practice

week of May 4-5: quiz + review of material since the prelim

**Final** (cumulative): Saturday May 9th at 9am.

**Thursday May 7:** 4820 review led by a few TAs *time tba*

- Friday, May 1<sup>st</sup>: divide and conquer
- Monday, May 4<sup>th</sup> ??

*Wed: SAT is NP-complete*

- CS theory club:  
<https://theoryclub.cs.cornell.edu/>
- **Wednesdays at 5:00 PM — CIS 450**
- **5:00-5:30 PM:** Reception/Social
- **5:30-6:30 PM:** Talk
- Pizza provided!
- Talk this Wednesday;: **Joey Rivkin**, first-year PhD student in theoretical computer science.
- **Defying Gravity with Range Minimum Queries**

The Range Minimum Query problem asks: given an array of integers, how do we answer the query: “What is the minimum element between indices  $i$  and  $j$ ?” We will see some beautiful ideas and data structures for answering many such queries efficiently, with runtimes that seem to defy intuition **gravity**.

# Today: Turing Machine: the simplest model of what is computation. Alan Turing (1936)

Recall: finite automata

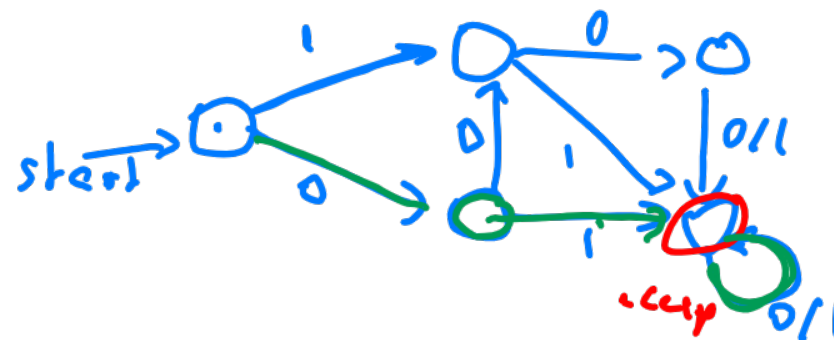
Graph

status:  $S$

start  $s \in S$

accept  $\epsilon \in S$

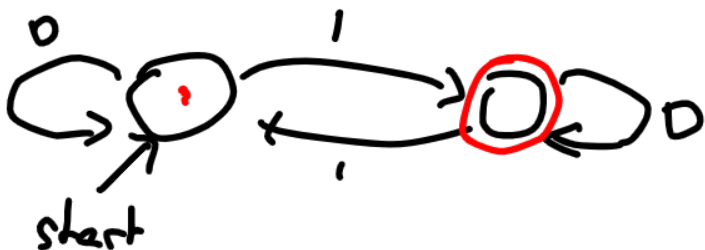
transition:  $S \times \sum \rightarrow S$   
 ↑  
 character



input 0110

edges labeled by characters

accept if end on a node labeled accept



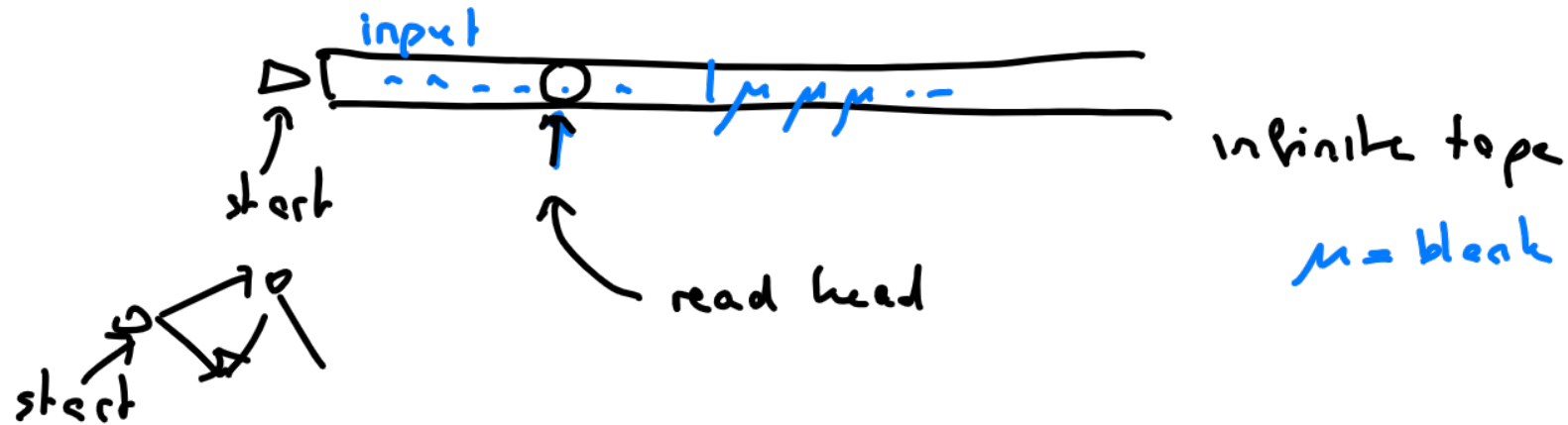
$\alpha \Rightarrow 00\dots010\dots0$

= strings with odd # of 1's

Recall 2800  $\alpha = \{0^n 1^n : n \text{ integer}\}$  00001111

# Today: Turing Machine: the simplest model of what is computation. Alan Turing (1936)

= finite automata with a tape



input :  $\Sigma$

$$\delta: S \times \Sigma \rightarrow S \times \Sigma \times \{+, 0, -\}$$

in state  $q \in S$  reading  $\sigma \in \Sigma$

change to new state  $q'$ , write  $\sigma'$  & move one character left/right/no move

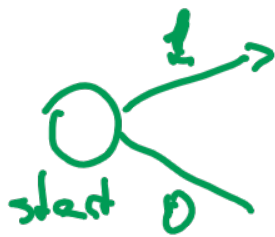
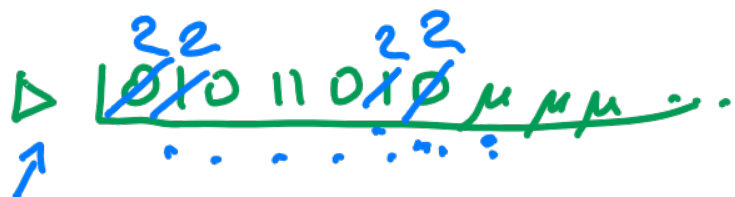
reach any accept state = accept

# Turing Machine: examples

Recognizing palindrome:

01011010

$\Sigma = \{0, 1, 2\}$



$\odot$  reject

start: first symbol  $\neq$  change to 2

go to the end

if last not the

same  $\Rightarrow$  not palindrome

if same  $\neq$  change to 2

same backwards on next  
letters ...

time:  $O(n^2)$

# Turing Machine: examples

$$L = \{0^n 1^n : n \text{ is even}\}$$



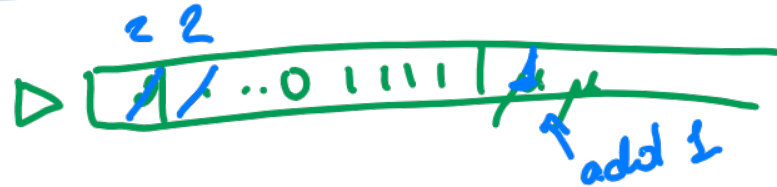
# steps:  $O(n^2)$

1. check is it  $0^n 1^m$

2. change one  $0 \rightarrow 2$  &  $1 \rightarrow 2$

accept if end no 0 or 1 left

Alternate:



$O(n^2)$

increment counter while seeing 0's

decrement —||— —||— 1

accept if counter ends at 0

# Church-Turing Hypothesis

$\Rightarrow$  anything computable  
is computable by a Turing machine

Poly time Python/Java =? poly time on Turing machine

? Quantum computers

quantum computable = computable by Python  
run time ? exponential

trouble: ① discrete log poly time on quantum computer  
given prime  $q$

② RSA

given two prime  $p+q$  (secret)

$n = pq$  - given: find  $p$  &  $q$

base  $g$

&  $A$

find  $a$ :  $g^a \equiv A \pmod{q}$

# Turing Machine more examples

$$\alpha = \{1^{2^k} : k \text{ integer}\}$$

$$\{1, 11, 1111, 11111111, \dots\}$$

#1's a power of 2

$$\Sigma = \{0, 1\}$$



1. check if input  $1^m$ 
  - if only one 1 ✓ accept
  - else check if #1's even  
reject if odd

change every other 1's to a 0

$\Rightarrow 2^{k-1}$  1s should be left

repeat: